

RAIL LINE ABANDONMENTS:  
NUMBER, LOCATION, AND TIME PHASING

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The U. S. Railway Association has recommended the closure of "potentially excess" rail lines in the Northeast and Midwest. Consequently, many state governments and local communities are in the process of evaluating which lines are worth subsidizing to preserve rail service. The abandonment of excess rail lines may benefit the rail industry by reducing some network costs such as rail maintenance and freight scheduling. Some of the cost savings from rail abandonment may be passed on to users in terms of lower freight rates and/or improved rail service. On the other hand, the abandonment of some lines is certain to impose additional economic and social costs on many rural areas.

As alternative rural transportation systems are evaluated the questions arise: Which of the existing rail lines should be abandoned; and, when should the lines be closed?

This paper presents a method of analysis to determine the number, location and time phasing of rail abandonments to minimize the discounted transportation costs of a community over a planning horizon of  $T$  years. A small scale version and numerical example are presented of a Stollsteimer-type location model used to determine the optimal locational pattern of rail lines and grain elevators in a region surrounding Fort Dodge, Iowa [2]. In contrast to the Iowa study, this paper focuses primarily on the questions of rail line abandonment and extends the Iowa study to determine the optimal location and time phasing of rail closures.

Number and Location of Rail Abandonments

Suppose there are  $G$  origins within a region. Each origin ships freight to one of  $M$  markets. Freight may be shipped directly from origin to market by truck; or, freight may be transshipped to market through a rail reload station. Freight that is transshipped to market is assembled to a rail station by truck and shipped from rail stations to market by rail. The supply of freight shipped from each origin is known; the supply received at each market is variable.

At the beginning of the planning horizon,  $R$  rail lines exist in the region. Some rail lines should, perhaps, be abandoned and others should be maintained.

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Transportation costs include a fixed annual maintenance cost required to ensure continuation of local rail services plus a variable cost of shipping by the least costly mode of transportation. The annual maintenance cost of rail lines varies by the number and location of rail lines under consideration and is borne by the community in the form of a subsidy.

Define the symbols

$\epsilon$  = element of

$S_g$  = origin or source  $g$ ;  $g = 1, 2, \dots, G$

$R_i$  = location of rail station  $i$ ;  $i = 1, 2, \dots, I$

$L_r$  = location of rail line  $r$ ;  $r = 1, 2, \dots, R$

$M_j$  = market  $j$ ;  $j = 1, 2, \dots, M$

$\lambda_{km}$  =  $k^{\text{th}}$  set of  $m$  rail lines,  $m = 0, 1, 2, \dots, R$ ; where  $m = 0$  denotes no rail lines.  $k = 1, 2, \dots, K_m \cdot K_m = R!/m!(R-m)!$

Note:  $\lambda$  is used as a shorthand notation for  $\lambda_{km}$

$X(g..)$  = quantity of freight shipped from  $S_g$

$X(gij)$  = quantity of freight transshipped from  $S_g$  to  $M_j$  through  $R_i$

$X(g.j)$  = quantity of freight shipped directly from  $S_g$  to  $M_j$

$X(..j)$  = quantity of freight shipped to  $M_j$

$$= \sum_{gi} X(gij) + \sum_g X(g.j)$$

$X(gi.)$  = quantity of freight shipped from  $S_g$  to  $R_i$

$C(gi.)$  = marginal cost of transporting freight from  $S_g$  to  $R_i$

$C(.ij)$  = marginal cost of transporting freight from  $R_i$  to  $M_j$

$C(gij) = C(gi.) + C(.ij)$

$\alpha(r)$  = annual equivalent cost of maintaining rail line  $L_r$

$TTC|\lambda$  = total truck plus rail transportation cost given rail line locational pattern  $\lambda_{mm}$

$$= \sum_r \alpha(r) + \sum_{gij} C(gij)X(gij) + \sum_{gj} C(g.j)X(g.j); r, i \in \lambda$$

$\beta(i)$  = per unit transfer (reload) cost of using rail station  $R_i$

$X(g..)$ ,  $X$ ,  $C(gi.)$ ,  $C(.ij)$ ,  $\alpha(i)$  and  $\beta(i)$  are known constants.  $X(g.j)$ ,  $X(gij)$ , and  $X(..j)$  are variables.

The problem to be solved is: Determine (a)  $m$ , the number of rail lines; (b)  $\lambda_{km}$ , the locational pattern of rail lines; and (c)  $X(g,j)$  and  $X(gij)$  the spatial flow of freight from origins to market to minimize total costs (TC) where

$$(1) \text{ TC} = \sum_{r \in \lambda} \alpha(r) + \sum_{g \in \lambda} \sum_j [C(gij) + \beta(i)] X(gij) \\ + \sum_{g,j} C(g,j) X(g,j)$$

This problem may be described as a transshipment rail location problem. The method of solution uses combinatorial search procedures initially developed by Stollsteimer [3] and later expanded by Ladd and Lifferth [2].

The problem is solved in two parts. The first part determines the spatial flow (allocation) of freight from origins to market for each  $\lambda_{km}$ . That is, part one is carried out once for each locational pattern of rail lines. Finding the optimal path of freight from each origin for a given network of rail lines is basically a problem of finding the shortest route in an acyclic network and, as such, may be determined using a recursive (sequential) search algorithm [4, p. 235]. Network  $\lambda_{km}$  is acyclic because freight is shipped from an origin to a rail station or to a market but not from origin to origin; and freight is shipped from a rail station to a market, but not from a rail station to an origin or to another rail station to be transshipped again.

Once the shortest route for each origin is found, the total variable transportation and reload costs for each locational pattern of rail lines, denoted as  $\text{TVC}|\lambda_{km}$ , is determined.

The second part determines the one locational pattern of rail lines for which total variable transportation and reload costs and fixed costs of rail line maintenance are minimized. The optimal number and location of rail lines are determined by systematically comparing total cost for each  $\lambda_{km}$  and selecting that constellation of rail lines for which joint total cost is minimum.

Part One determines  $\text{TVC}|\lambda_{km}$  as follows: Let  $\text{AVC}(g_i g_j)$  denote the average variable transportation and reload cost of shipping one unit of freight over the shortest route from origin  $g$  to market if freight is transshipped through a rail station. For  $R_{i \in \lambda}$  compute

$$(2) \text{ AVC}(g_i g_j) = \min_i \{ \min_j [C(i,j) + \beta(i)] + C(g_i) \} \text{ for all } g. \text{ Expression}$$

(2) provides a  $G \times 1$  matrix of  $[\text{AVC}(g_i g_j)]$

Denote  $\text{AVC}(g, j_g)$  as the lowest attainable average variable cost of shipping freight from origin  $g$  to market if freight is transported directly from origin to market. Compute  $\text{AVC}(g, j_g)$  for all  $g$  as follows:

$$(3) \text{ AVC}(g, j_g) = \min_j C(g, j)$$

Let  $\text{AVC}(g, \dots)|\lambda$  denote the shortest route from origin  $g$  through network  $\lambda_{km}$  and compute as follows:

$$(4) \text{AVC}(g..)|_{\lambda_{km}} = \min [\text{AVC}(g_i.j_i), \text{AVC}(g.j_g)]$$

Total variable cost given  $\lambda_{km}$  is

$$(5) \text{TVC}|_{\lambda_{km}} = \sum_g \text{AVC}(g..) X(g..)|_{\lambda_{km}}$$

Part Two determines the locational pattern of rail lines and spatial flow of freight that minimizes (1). Let  $\overline{\text{TC}}|_{\lambda_{km}}$  denote the minimum total cost given  $\lambda_{km}$  and compute as follows:

$$(6) \overline{\text{TC}}|_{\lambda_{km}} = \text{TVC}|_{\lambda_{km}} + \sum_{r \in \lambda_{km}} \alpha(r)$$

Compute (6) for all  $k$  and for all  $m$ . The minimum total cost from  $m$  rail lines is

$$(7) \overline{\text{TC}}|_m = \min_k \overline{\text{TC}}|_{\lambda_{km}}$$

The number and locational pattern of rail lines that minimizes total cost is

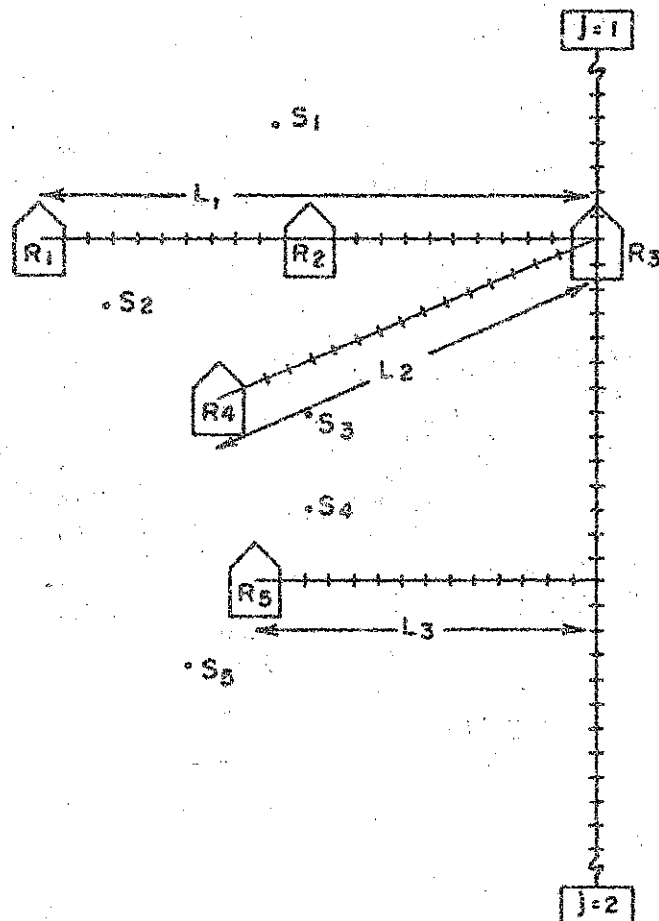
$$(8) \overline{\text{TC}} = \min_m \overline{\text{TC}}|_m$$

### Example 1

The method of solution used to solve the rail location problem can be illustrated as follows: Suppose there are five origins, two markets, five rail stations, and three rail lines classified as "potentially excess" that connect with a major trunk line. The supply of freight at origins is known. Freight may be shipped from origins to market directly, or it may be trans-shipped to market through rail stations.

Figure 1 shows the location of origins, rail stations, final markets, and potential rail linesites. Final destinations are denoted by  $\square$ ; rail stations are denoted by  $\triangle$ ; dots,  $\cdot$ , represent origins; and +++++ represents a rail line.

FIGURE 1. LOCATIONAL PATTERN OF ORIGINS, RAIL STATIONS, MARKETS, AND RAIL LINES FOR  $\lambda_{13}$ .



Data used for this example are presented in the following order: (a) alternative rail locational patterns; (b) supply at origins; (c) transportation costs; and (d) reload costs at rail stations.

Rail locational options: alternative rail line locational patterns used in this example are defined as follows:

- $\lambda_{10}$  = no feeder lines
- $\lambda_{11}$  = rail line  $L_1$
- $\lambda_{21}$  = rail line  $L_2$
- $\lambda_{31}$  = rail line  $L_3$
- $\lambda_{12}$  = rail lines  $L_1$  and  $L_2$

$\lambda_{22}$  = rail lines  $L_1$  and  $L_3$

$\lambda_{32}$  = rail lines  $L_2$  and  $L_3$

$\lambda_{13}$  = rail lines  $L_1$ ,  $L_2$  and  $L_3$

The rail line network defined by  $\lambda_{13}$  is presented in Figure 1.

Supply at origins: The volume of freight shipped from each origin is predetermined.

$X(1..) = 100,000$ ;  $X(2..) = 100,000$ ;  $X(3..) = 100,000$ ;  $X(4..) = 100,000$ ;  
 $X(5..) = 100,000$

Transportation costs: the per unit transportation cost of shipping from origin  $S_g$  to rail station  $R_i$ , and from  $S_g$  to  $M_j$  are presented in Table 1. Table 2 presents the per unit transportation costs of shipping freight from rail station  $R_i$  to market  $M_j$ . The annual cost of maintaining rail line  $L_1$  is  $\alpha(1) = \$275,000$ ;  $L_2$  is  $\alpha(2) = \$325,000$ ; and  $L_3$  is  $\alpha(3) = \$125,000$ .

Table 1. TRANSPORTATION COSTS FROM ORIGINS TO RAIL STATIONS, AND FROM ORIGINS TO MARKETS IN DOLLARS PER UNIT:  $C(g_i.)$  and  $C(g.j)$

Source	Rail station					Market	
	$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$M_1$	$M_2$
$S_1$	1.0	.4	1.2	1.0	1.6	5.8	7.4
$S_2$	.3	.8	1.8	.6	1.2	4.4	7.0
$S_3$	1.2	.6	1.2	.3	.6	6.8	6.8
$S_4$	1.4	1.0	1.5	.5	.3	7.2	6.6
$S_5$	1.5	1.4	2.0	.8	.2	7.6	6.0

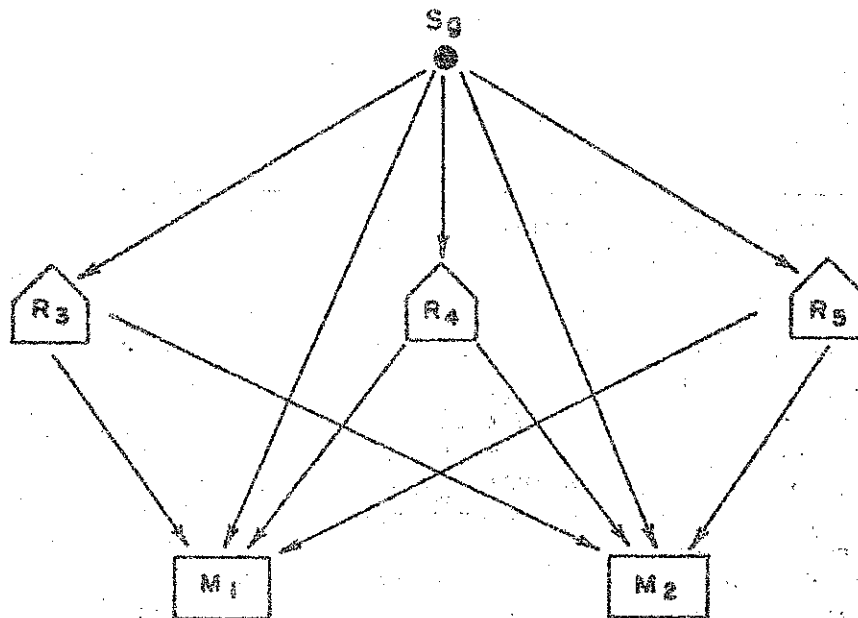
Table 2. AVERAGE VARIABLE TRANSPORTATION COSTS FROM RAIL STATIONS TO MARKET IN DOLLARS PER UNIT,  $C(.ij)$

Rail station	Market	
	$M_1$	$M_2$
$R_1$	2.0	2.3
$R_2$	1.8	2.1
$R_3$	1.5	1.8
$R_4$	1.8	2.1
$R_5$	2.1	1.8

Reload costs: The reload cost at rail stations represents the per unit cost of transferring freight from one mode to rail. For this example  $\beta(i) = \$0.50$  for all  $i$ .

Part One determines the optimal routing of freight from origin  $g$  and the total variable cost for each constellation of rail lines. The location of rail stations is also determined in part one and depends upon the number and location of rail lines. The routing options, for example, of  $S_g$  given  $\lambda_{32}$  are illustrated schematically in Figure 2. For  $\lambda_{32}$  freight may be transshipped from  $S_g$  through  $R_3$ ,  $R_4$ , or  $R_5$ . For  $\lambda_{11}$  freight may be transshipped from  $S_g$  through  $R_1$ ,  $R_2$ , or  $R_3$ .

FIGURE 2. SHIPPING OPTIONS AVAILABLE FROM ORIGIN  $S_g$  GIVEN  $\lambda_{32}$



The optimal routing pattern for each  $\lambda_{km}$  is determined using a sequential routing algorithm. The first step is to determine the best route from rail station  $R_i$  to market for all  $i$ . Table 3 presents an  $I \times J$  matrix of  $AVC(.ij)$  where  $AVC(.ij) = C(.ij) + \beta(i)$ . The best market for  $R_i$  is found by  $\min [C(.ij) + \beta(i)] = AVC(.ij_1^*)$ . The symbol  $*$  is used to denote  $AVC(.ij_1^*)$  in the matrix  $AVC(.ij)$ . For example, since  $[AVC(.1,1) = 2.5]$  is less than  $[AVC(.1,2) = 2.8]$ ,  $M_1$  is the best market for  $R_1$  and all freight received at  $R_1$  should be shipped to  $M_1$  to minimize costs.



Table 3. AVERAGE VARIABLE COST OF RELOADING FREIGHT  
 AT  $R_i$  SHIPPING FREIGHT FROM RAIL  
 STATION  $R_i$  TO MARKET  $j$  WHERE  
 $AVC(.ij) = C(.ij) + \beta(i)$

Rail station	Market	
	$M_1$	$M_2$
$R_1$	2.5*	2.8
$R_2$	2.3*	2.6
$R_3$	2.0*	2.3
$R_4$	2.3*	2.6
$R_5$	2.6	2.3*

\* Denotes  $AVC(.ij_i) = \min_j AVC(.ij)$

The next step computes a matrix of  $[AVC(gij_i) = AVC(.ij_i) + C(gi.)]$  for each rail locational pattern  $\lambda_{km}$ . Table 4 presents  $AVC(gij_i)$  for the three possible combinations of one rail line; i.e.  $\lambda_{1,1}$ ,  $\lambda_{2,1}$ , and  $\lambda_{3,1}$ . The best route for  $S_g$  when freight is transshipped through a rail station to market is denoted by \* where  $AVC(gi_j_i) = \min_i AVC(gij_i)$ . For example,  $AVC(1,1,j_1) | \lambda_{1,1}$  represents the average variable cost of shipping freight from origin  $S_1$  to rail station  $R_1$  plus the transfer costs at  $R_1$  plus the average variable cost of shipping from  $R_1$  to the best market. From Table 3 it was determined that  $M_1$  is the best market for  $R_1$ . Thus,  $AVC(1,1,j_1) | \lambda_{1,1} = AVC(.1,1_1) | \lambda_{1,1} + C(1,1.) = 2.5 + 1.0 = 3.5$ . Repeating this step for each  $\lambda_{km}$  completes the operations of expression (2).

Table 4. AVERAGE VARIABLE COST FROM TRANSSHIPPING  
 FREIGHT FROM ORIGIN  $S_g$  THROUGH RELOAD  
 STATION  $R_i$  TO MARKET  $M_j$  FOR ALL  
 COMBINATIONS OF ONE ( $M = 1$ ) RAIL LINE; WHERE  
 $AVC(gij_i) | \lambda_{km} = AVC(.ij_i) + C(gi.)$

Source	$\lambda_{11}$			$\lambda_{21}$		$\lambda_{31}$	
	$R_1$	$R_2$	$R_3$	$R_3$	$R_4$	$R_3$	$R_5$
$S_1$	3.5	2.7*	3.2	3.2*	3.3	3.2*	3.9
$S_2$	2.8*	3.1	3.8	3.8	2.9*	3.8*	3.9
$S_3$	3.7	3.1*	3.2	3.2	2.6*	3.2*	3.2
$S_4$	3.9	3.3*	3.5	3.5	2.8*	3.5	2.9*
$S_5$	4.0	3.7*	4.0	4.0	3.1*	4.0	2.8*

\* Denotes  $AVC(gi_j_i) | \lambda_{km} = \min_i AVC(gij_i) | \lambda_{km}$

The values for expression (3) are found in Table 5. Table 5 is a  $G \times J$  matrix of  $[AVC(g,j) = C(g,j)]$ . The best route for  $S_g$  when freight is shipped directly to market is denoted by \* where  $AVC(g,j_g) = \min_j C(g,j)$ .

Table 5. AVERAGE VARIABLE COST OF SHIPPING FREIGHT DIRECTLY FROM SOURCE  $S_g$  TO MARKET  $M_j$ , WHERE  $AVC(g,j) = C(g,j)$

Source	Market	
	$M_1$	$M_2$
$S_1$	5.8*	7.4
$S_2$	4.4*	7.0
$S_3$	6.8*	6.8
$S_4$	7.2	6.6*
$S_5$	7.6	6.0*

\*Denotes  $AVC(g,j_g) = \min_j C(g,j)$

Expression (4) is found by selecting, for each  $\lambda_{km}$ , the best route for origin  $S_g$  given  $AVC(g_i, j_i)$  and  $AVC(g, j_g)$ . Table 6 presents the minimum attainable average variable cost for each  $S_g$  given  $\lambda_{km}$ .

Table 6. THE LOWEST ATTAINABLE AVERAGE VARIABLE COST OF SHIPPING FREIGHT FROM ORIGIN  $S_g$  FOR SELECTED RAIL LINE LOCATIONAL PATTERNS,  $AVC(g..) | \lambda_{km}^a$

Source	Rail line locational options							
	$\lambda_{10}$	$\lambda_{11}$	$\lambda_{21}$	$\lambda_{31}$	$\lambda_{12}$	$\lambda_{22}$	$\lambda_{32}$	$\lambda_{13}$
$S_1$	3.2	2.7	3.2	3.2	2.7	2.7	3.2	2.7
$S_2$	3.8	2.8	2.9	3.8	2.8	2.8	2.9	2.8
$S_3$	3.2	3.1	2.6	3.2	2.6	3.1	2.6	2.6
$S_4$	3.5	3.3	2.8	2.9	2.8	2.9	2.8	2.8
$S_5$	4.0	3.7	3.1	2.8	3.1	2.8	2.8	2.8

<sup>a</sup> See expression (4).

Table 7 presents total variable cost, fixed cost of rail maintenance, and the minimum attainable total cost for each rail line locational pattern. Total variable cost,  $TVC | \lambda_{km}$ , is found by computing  $\sum_g AVC(g..) X(g..) | \lambda_{km}$ .

Table 6 presents  $AVC(g..)$  for each  $\lambda_{km}$ ; and, the volume of freight at  $S_g$ ,  $X(g..)$ , is predetermined.

Table 7. TOTAL VARIABLE COST, FIXED COST OF RAIL LINE MAINTENANCE,  
AND MINIMUM ATTAINABLE TOTAL COST FOR SELECTED RAIL  
LOCATIONAL PATTERNS, IN THOUSANDS OF DOLLARS<sup>a</sup>

	Rail line locational options							
	$\lambda_{10}$	$\lambda_{11}$	$\lambda_{21}$	$\lambda_{31}$	$\lambda_{12}$	$\lambda_{22}$	$\lambda_{32}$	$\lambda_{13}$
$TVC   \lambda_{km}$	3,540	3,120	2,920	3,180	2,800	2,860	2,860	2,740
$\sum_{r \in \lambda_{km}} \alpha(r)$	0	275	325	125	600	400	450	725
$\overline{TC}   \lambda_{km}$	3,540	3,395	3,245	3,305	3,400	3,260	3,310	3,465

<sup>a</sup> See expressions (5) and (6).

Part Two determines the optimal number and location of rail lines that minimizes total variable costs plus total fixed costs of rail maintenance. From Table 7 the minimum total cost from M rail lines can be identified. The number and location of rail lines that provide the lowest attainable cost for expression (1) is also identified from Table 7. Locational pattern  $\lambda_{21}$  with a total systems cost of \$3,245,000 represents the optimal number and location of rail lines.

It is not necessary to compare every possible combination of rail lines to find a global optimum. If the best location of 4 rail lines is better than the best location of both 3 rail lines and 5 rail lines, then the best location of 4 rail lines is better than any other number and location of rail lines. The change in total variable costs between the best location of M rail lines and the best location of M + 1 rail lines will never be positive. The best location of M rail lines is included in the different locational combinations of M + 1 rail lines. An increase in total variable costs from adding an additional rail line can always be avoided by selecting the best location of M rail lines and refusing to use the M + 1st rail line. Thus, once a local optimum is found, additional comparisons are unnecessary since any local optimum is also global.

In this simple illustration the trade-off between rail maintenance costs and shipping costs is demonstrated. As rail lines are abandoned there are some savings of rail maintenance costs. As rail lines are abandoned, however, freight must be shipped from origins by truck to more distant rail stations for transshipment to market; or, freight must be trucked directly from origin to market. The additional cost of shipping freight resulting from rail abandonment must be balanced against the cost savings of maintaining fewer lines.

### Time Phasing of Rail Abandonments

Expressions (1) through (8) describe a method of solving for the optimal number and location of rail lines for a given volume of freight and for a given level of costs. Changes in the parameters of the model, however, may influence both the pattern of shipments that minimizes total cost as well as the locational pattern of rail lines. Rail abandonment decisions, therefore, that are based only on current economic conditions may not be in the long run interest of the public if changes in the volume of freight and/or level of transportation costs are anticipated. Abandoning a rail line, in most cases, is an irreversible decision. On the other hand, rail abandonment decisions should not be based on future transportation needs disregarding current conditions. Maintaining excess rail lines, for future use, for example, may impose costs that are greater than the long run benefits.

Suppose, for example, that the data  $[X(g.), C(gi.), C(.ij), \alpha(r), \text{ and } \beta(i)]$ , used in example 1 reflect the economic conditions prevailing during period 1; and further, suppose that the volume of freight increases from the period 1 level of 200,000 units per origin to 500,000 units per origin in period 2. The best locational pattern for period 1, disregarding period 2, is  $\lambda_{21}$ . The best locational pattern for period 2, disregarding period 1, is  $\lambda_{22}$ . The problem is clear: If rail abandonment decisions are irreversible and if rail lines  $L_1$  and  $L_3$  are abandoned in period 1 to minimize total system cost for the users in period 1, then  $\lambda_{22}$  is not a feasible alternative for period 2. As an alternative,  $\lambda_{22}$  may be implemented during period 1 and maintained for use in period 2. In this case, the additional cost of maintaining  $\lambda_{22}$  during period 1 ( $TC|\lambda_{22} - TC|\lambda_{21} = 5,000$ ) must be balanced against the discounted value of  $\lambda_{22}$  over  $\lambda_{21}$  in period 2.

The procedures outlined below extend the rail location model presented in expressions (1) through (8) to account for the temporal as well as the spatial interdependencies that influence the number and location of rail lines. The problem is to determine (a) the number, location and time phasing of rail abandonments, and (b) the flow of freight from origins to market for each period to minimize discounted total costs over a finite time horizon of T years. The method of solution involves the following four steps:

Step 1: Compute the minimum total system cost in time t for each rail locational network  $\lambda_{km}$  given the predetermined level of costs and volume at origins during time t;  $t = 1, 2, \dots, T$ . Let  $TC(t)|\lambda_{km}$  denote the minimum total cost in time t given  $\lambda_{km}$ .  $TC(t)|\lambda_{km}$  is determined by solving expressions (1) through (6) for each period t. Predetermined data  $[X(g.), C(gi.), C(.ij), \alpha(r), \text{ and } \beta(i)]$  are changed for each t to reflect the economic conditions prevailing during time t.

Step 2: Construct a decision tree representing the feasible rail network options over a time horizon of T periods. At the start of time  $t = 0$  there is 1 decision node; at the end of time  $t = 1$  there are  $\sum_m^K$  possible decision nodes. Each node represents one rail line network as  $\lambda_{km}$  defined by  $\lambda_{km}$  where  $k = 1, 2, \dots, K_m$  and  $m = 0, 1, \dots, R$ . At the end of time  $t = 2$  there is another set of decision nodes (or rail network options) which branch from the nodes of the previous period. The number of decision nodes in time  $t = 2$  depends upon the feasible rail network options available in period 2 given the combinations of rail networks in period 1. This branching process continues until all feasible adjustment paths over time are identified.

Step 3: Label each decision node. Let  $N_{st}$  denote decision node  $s$  in time  $t$ ;  $s = 1, 2, \dots, S_t$ . During time  $t = 0$ ,  $S_t = 1$ . During time  $t = 1$ ,  $S_t = \sum_{m=1}^K K_m$ . During other periods,  $t > 2$ ,  $S_t$  depends upon the number of network decision nodes feasible during the period.

Step 4: Compute the discounted total cost of alternative rail abandonment options and select that adjustment path which minimizes discounted total cost.

Let

$$\overline{TC}|_{N_{st}} = \overline{TC}(t) | \lambda_{km} \text{ where}$$

$\lambda_{km}$  corresponds to  $N_{st}$ , by definition as specified in Step 3.

= minimum total cost in time  $t$  given the rail locational network identified by decision node  $N_{st}$ .

$D$  = single period discount rate.

$\overline{DTC}|_{N_{st}}$  = minimum discounted total cost at  $N_{st}$  from  $t$  to  $T$ .

$\overline{DTC}|_{N_{10}}$  represents the minimum discounted total cost over a finite time horizon of  $T$  periods.  $\overline{DTC}|_{N_{10}}$  is found by solving expression (9) for each time period beginning with the last period ( $t = T$ ) and working backwards (sequentially) through time.

$$(9) \overline{DTC}|_{N_{st}} = D \min_{s \in N_{st}} [\overline{DTC}|_{N_{s, t+1}}] + \overline{TC}|_{N_{st}}$$

for  $s = 1, 2, \dots, S_t$  and  $t = T - 1, T - 2, \dots, 1, 0$ .

where

$$\overline{DTC}|_{N_{st}} = \overline{TC}|_{N_{st}}$$

$$\text{and } \overline{TC}|_{N_{10}} = 0$$

### Example 2

Steps 1 through 4 are illustrated as follows. Let the data of example one represent the conditions prevailing during period one. Assume that in period two the volume at each origin increases from the period 1 level to:  $X(1..) = 500,000$ ;  $X(2..) = 500,000$ ;  $X(3..) = 500,000$ ;  $X(4..) = 500,000$ ; and  $X(5..) = 500,000$ . Table 8 presents  $\overline{TC}(t) | \lambda_{km}$  for  $t = 1$  and  $t = 2$ .

Table 8. MINIMUM ATTAINABLE TOTAL COST FOR TIME 1 AND TIME 2  
GIVEN SELECTED RAIL LOCATIONAL PATTERNS, IN  
THOUSANDS OF DOLLARS,  $TC(t) | \lambda_{km}$

	$\lambda_{10}$	$\lambda_{11}$	$\lambda_{21}$	$\lambda_{31}$	$\lambda_{12}$	$\lambda_{22}$	$\lambda_{32}$	$\lambda_{13}$
Time $t = 1$	3,540	3,395	3,245	3,305	3,400	3,260	3,310	3,465
Time $t = 2$	8,850	8,075	7,625	8,075	7,600	7,550	7,600	7,575

Step 2 identifies the various rail line abandonment options that are feasible over time. In Step 3 the various nodes of the decision tree identified in Step 2 are labeled. Figure 3 presents the decision tree of the feasible rail line abandonment options over two time periods for  $m = 0, 1,$  and 2. Decision nodes are identified by both  $\lambda_{km}$  and  $N_{st}$ .

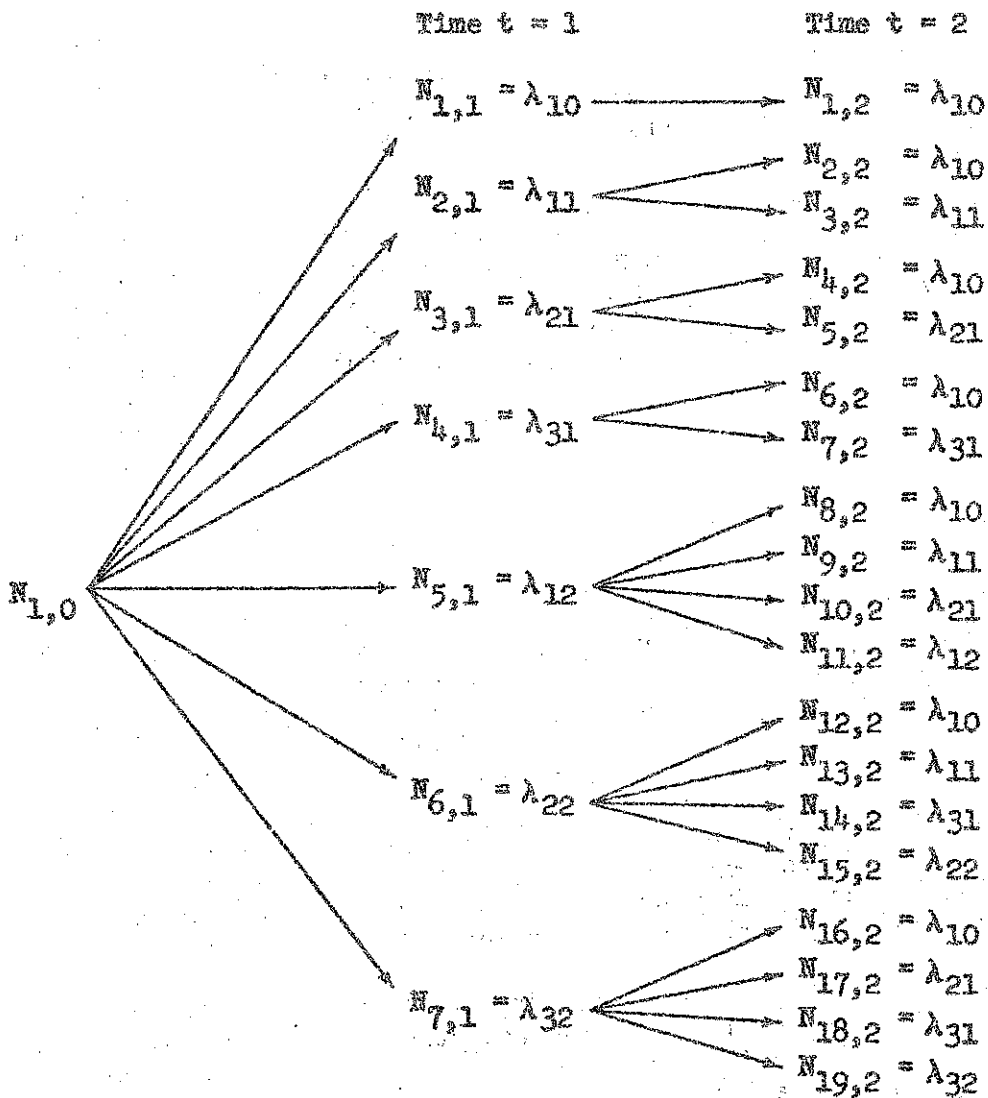


Figure 3. RAIL LINE ABANDONMENT OPTIONS OVER A TIME HORIZON OF 2 PERIODS

Step 4 identifies the optimal adjustment path of rail abandonments over time. The optimal time phasing of abandonment is determined by selecting that adjustment path that minimizes discounted total costs.  $DTC|N_{10}$  is found by sequentially computing  $DTC|N_{st}$  for each time period beginning with  $t = T$  and working backwards through time. For  $t = 2, = T$ ,  $DTC|N_{s,2}$  can be read directly from Table 8; e.g.,  $DTC|N_{1,2} = 8,850$ ;  $DTC|N_{2,2} = 8,850$ ;  $DTC|N_{3,2} = 8,075$ ; ...,  $DTC|N_{19,2} = 7,600$ . For  $t = 1$ ,  $DTC|N_{s1}$  is computed as described by expression (9). For example,  $DTC|N_{2,1} = \{D \min [DTC|N_{2,2}; DTC|N_{3,2}] + TC|N_{2,1}\} = \{.9091 \min [8,850; 8,075] + 3,395\} = 10,736$  where  $D = 1/(1 + .10)$ .

This step is repeated for each  $N_{s1}$ ;  $s = 1, 2, \dots, S_1$ . Step 4 is completed when  $DTC|N_{1,0}$  is determined.

$$\begin{aligned} DTC|N_{1,0} &= \{D \min [DTC|N_{1,1}; DTC|N_{2,1}; \dots; DTC|N_{6,1}; \dots; DTC|N_{8,1}] + 0\} \\ &= \{.9091 \min [11,586; 10,736; \dots; 10,124; \dots; 10,329] + 0\} \\ &= .9091 (10,124) = \$9,204 \end{aligned}$$

In this example, discounted total cost was minimized by selecting  $DTC|N_{6,1}$  which represents implementing  $\lambda_{22}$  during period 1 and  $\lambda_{22}$  during period 2. Any other number, location, and time phasing of rail abandonments would be suboptimal.

In summary, the number and location of rail lines to minimize total transportation costs during period one, disregarding period two, is defined by option  $\lambda_{21}$  with one rail line located at  $L_2$ . When the future conditions of period two are considered along with period one in making the rail abandonment decisions, option  $\lambda_{22}$  is best. Locational pattern  $\lambda_{22}$  includes rail lines  $L_1$  and  $L_3$ . The additional cost of maintaining too many lines during period one is less than the additional cost of having too few lines in period two.

#### Application and Conclusion

A variation of the model presented in this paper was used to determine the optimal locational pattern of rail lines and grain elevators in a region surrounding Fort Dodge, Iowa. Joint net revenue of grain producers was maximized by abandoning all feeder (or branch) rail lines in the region and hauling freight by truck from origins to subterminal elevators located on major trunk lines. The cost savings of maintaining fewer rail lines plus the quantity discounts available for shipping larger volumes of grain from subterminal elevators were balanced against the additional trucking costs plus the investments required to establish subterminals on trunk rail lines. Additional energy requirements, road use costs, and pollution emission resulting from rail line abandonment were also estimated.

The planning horizon extended from 1971 to 1980. The location and capacity of grain handling and transportation facilities existing at the beginning of the planning horizon were taken into account in designing the transportation network. In the Fort Dodge study the problem of estimating

the optimal phasing of rail abandonments over the planning horizon was a moot question. The best locational pattern of rail lines in 1980, disregarding other time periods, was identical to the best locational pattern of rail lines in 1971, given the level of costs and volume of freight existing in 1971. Discounted joint net revenue of producers was maximized, therefore, by abandoning all feeder (branch) rail lines at the beginning of the planning horizon. Detailed results of the study are reported in Baumel, et al. [1].

Many communities or geographical regions in the Northeast and Midwest are presently sustaining rail lines that have been classified by the U.S. Railway Association as "potentially excess." The classification of such lines, however, has been based primarily on single line analysis without regard to intramodal interdependencies that exist over space as well as over time. The procedures for multiple-line analysis presented in this paper provide one method of including intra- and inter-modal interdependencies in the evaluation and design of alternative rural transportation systems.



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